

# Patterns in Combinator Evolution

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Rules for speeding up SK combinator evaluation were investigated, and experiments were performed to determine the proportion of SK combinator expressions that reach a fixed point before exceeding certain computational limits. It was found that approximately 80.3 percent of SK combinator expressions with size 100 reach a fixed point within 300 steps and also without having reached a combinator size of 200 000 sometime during evaluation.

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## 1. Introduction

Combinators are constructs that can be understood as functions in a certain abstract sense [1] and that are useful as a simplified model of computation. There are two very interesting combinators called S and K, which can be used together to write any computer program (the two together are Turing complete, so can be used as a universal language).

To work with these combinators, we can look at transformations of combinator expressions. This involves starting out with an initial combinator expression, which is made by calling S and K combinators on each other in some order. Once an SK combinator expression has been generated, we apply the rules for evaluating the combinators repeatedly. The rules are to replace  $K[x_][y_]$  with  $y$ , and  $S[x_][y_][z_]$  with  $x[z][y[z]]$ . So for example, we start out with the expression  $S[S[S]][S][K][K]$ , and running it step by step, we would get

```
S[S[S]][S][K][K]
S[S][K][S[K]][K]
S[S[K]][K][S[K]][K]
S[K][K][K][S[K]][K]
K[K][S[K]][K]][K][K][S[K]][K]]
K[S[K]][K]
S[K]
S[K]
```

The evaluation “halts” when applying the rules no longer changes the expression, so the “output” of  $S[S[S]][S][K][K]$  is  $S[K]$ .

One detail that was glossed over in the previous explanation is the order in which rules are applied to the expression. There are two standard ways in which to apply the rules, normal order and applicative

order. In applicative order, the rules are applied at each step by scanning once from left to right and applying the rule wherever possible without overlapping. This leads to larger parts of the combinator having the rules applied to them first. In normal order, the rules are again applied from left to right without overlapping, but the parts deepest inside brackets (the arguments of the functions) are evaluated before the outer parts.

Normal order will be used exclusively in this paper. It turns out that there are often combinator expressions that reach a fixed point in normal order but not in applicative order. This is because sometimes there will be a subpart of a combinator that would continuously expand if evaluated over and over, but if that subpart is inside a K combinator, evaluating the K function before the arguments deletes the whole expression that was going to grow before it could begin its continuous growth. For example,  $S[S[S]][S][S][S][S]$  never stops growing [2], so evaluating  $K[S][S[S]][S][S][S][S]$  in applicative order will never stop growing either, whereas evaluating the same in normal order will cause the expression to become just S and to then stop changing, after just one step. In this case, the K simply deletes the second part containing the potential to grow.

In general, some combinator expressions reach a fixed point and stop changing, which is referred to as “halting,” whereas some never stop growing, or begin to continuously loop through a set of states.

This paper discusses how certain patterns in combinator evaluation were discovered. Extra combinator rules based on these patterns were used in attempts to speed up the evaluation of large combinators. Also, experiments were performed to determine the proportion of SK combinator expressions that reach a fixed point before exceeding certain computational limits.

In [3], the authors created databases of combinator reductions to speed up computation. Here a slightly different approach to finding extra combinator rules for speeding up combinator reductions is used. In [3], a sort of “multiplication table” was built up of different ways combinators can evaluate, whereas in this paper, combinators up to a certain size were evaluated with different combinator parts being replaced by variables, as described in Section 2.

## 2. Patterns in Combinator Evolution

Out of all possible SK combinator patterns up to size five (of which there were 2582), just 13 are non-redundant as defined in a certain way. These 13 rules are rules in which evaluating the left side to get the right side takes more than just one step. Also, less general rules

are excluded in favor of more general rules (for instance, the more general  $S[K][y\_][x\_]\rightarrow x$  is included but not the more specific  $S[K][S][x\_]\rightarrow x$ ). Also, cases where applying a shorter combinator rule has the same effect as applying a larger combinator rule are not included. The rule  $S[S[K][y\_][x\_]]\rightarrow S[x]$  is not included, since the shorter rule  $S[K][y\_][x\_]\rightarrow x$  applied to  $S[S[K][y\_][x\_]]$  will also give  $S[x]$  in one step. The 13 non-redundant rules, which are general ways in which combinators up to size five evaluate, are the following:

```

K[x_][y_] → x
S[x_][y_][z_] → x[z][y[z]]
S[K][x_][y_] → y
K[K[x_][y_]][z_] → x
K[K[x_]] [y_][z_] → x
K[K][x_][y_][z_] → y
S[x_][K[y_]][z_] → x[z][y]
S[K[x_]] [y_][z_] → x[y[z]]
S[S][x_][y_][z_] → y[z][x[y][z]]
S[S[x_]] [y_][z_] → x[y[z]][z[y[z]]]
S[S][x_][K[y_]] → y
S[S[K]] [x_][y_] → x[y]
S[S][K][x_][y_] → x[y][x]

```

There are 22 994 possible rules of size six, and again a relatively small number, 67, are non-redundant in this way.

### 3. Applying Optimization

Next there is the matter of applying the rules that were found to combinator expressions. Combinators of various sizes were randomly generated. These combinators were then evaluated, while adding extra rules to see if this sped up the combinator evolution. It is expected that the more steps and the more often rules are applied, the more extra rules would help speed up evolution. Large combinators generally have more rule applications and run for more steps. Larger and larger combinators were tested to look for any speedups.

There is one issue with using large combinators, however, which is discussed more in Section 4. Basically, the issue is that larger combinators were more likely not to reach a fixed point within 300 steps. Generally, only combinators that reach a fixed point were of interest for testing optimizing rules, since the extra rules would likely be used primarily for evaluating halting combinators. (If we were looking at nonhalting combinators, we would presumably want to look at step-by-step evolution, and the extra rules would likely be skipping steps in unpredictable increments.)

Based on some preliminary experimenting into how large combinators evolved, combinators that seemed unlikely to halt in two ways were first analyzed. Combinators that did not halt after a certain num-

ber of steps and combinators that did not halt before growing to a certain size were selected. As elaborated more in Section 4, larger combinators do seem to either quickly begin to grow exponentially or to quickly begin evolving in a repetitive way. A small minority of combinators that did neither of these things, but instead grew non-exponentially, but in a complex way for many steps, were ignored. (These are rather interesting combinators; however, it seems especially hard to predict what they will do in the future: whether they will continue having complicated slow growing behavior, whether they will begin growing exponentially at some point, or whether they will in the end exhibit repetitive behavior after all.)

However, even with combinators of size 1000, the extra rules only modestly sped up evaluation. With only S and K rules, Mathematica took an average of 1.20 seconds each to repeatedly evaluate 200 randomly generated size 1000 combinators (these were also selected as combinators that reached a fixed point before reaching either 10 000 iterations or a size of 10 000). With the extra 11 rules, this took 1.13 seconds, and with only one extra rule,  $S[K][x\_][y\_]\rightarrow y$ , this took 1.06 seconds. All timing measurements here were performed with Mathematica 11 on a 3.40 GHz Intel Core i7-6700 processor.

Different selections and orders of the 11 optimization rules found, along with the original S and K rules, were tested to see which would speed up evaluation the most. First, each extra rule was added one at a time, and the added rule was added at almost each possible index in the rule list. Some positions were not checked. For example, putting the extra rule  $S[K][x\_][y\_]\rightarrow y$  after  $S[x\_][y\_][z\_]\rightarrow x[z][y[z]]$ , as opposed to before it, would mean the extra rule would not be used. The extra rule would be skipped over by the original and more general rule. The fastest rule permutation turned out to be  $\{K[x\_][y\_]\rightarrow x, S[K][x\_][\_]\rightarrow x, S[x\_][y\_][z\_]\rightarrow x[z][y[z]]\}$ , which took an average of 1.00 second to evaluate the same combinators as before.

This is interesting: since  $S[K][x\_][y\_]\rightarrow y$  is equivalent to the identity or “I” combinator,  $(I[y\_]\rightarrow y, I$  is essentially the same as  $S[K][x\_]$ ), which can be used in practice to simplify the use of SK combinators. Extra rules were added in all possible combinations. Two, three and then four different rules were added. These are the fastest rules in each of those situations, with rules on the left and time in seconds on the right of each list:

```
RepeatedTiming[a /. {b -> c}]
{9.8 × 10-7, a}
RepeatedTiming[a /. {b -> c, d -> e}]
{1.4 × 10-6, a}
RepeatedTiming[a /. {b -> c, d -> e, e -> f}]
{2.02 × 10-6, a}
```

The extra rules can be used for speeding up combinator evaluation. However, there is a tradeoff when using them, as more rules can help

in skipping steps, but extra rules also generally slow down evaluation. This is due to the interpreter needing to scan for more rule matches.

With only S and K rules, Mathematica took an average of 1.20 seconds each to repeatedly evaluate 200 randomly generated size 1000 combinators that were also selected as combinators that reached a fixed point before reaching either 10 000 iterations or a size of 10 000. With the extra 11 rules, this took 1.13 seconds, and with only one extra rule,  $S[K][x\_][y\_]\rightarrow y$ , this took 1.06 seconds.

When only one to four extra rules were used, the rule selection and order that were fastest for each were found to be the following:

```
extra1 → {K[x_][y_] → x, S[K][x_][y_] → y, S[x_][y_][z_] → x[z][y[z]], 1.}
extra2 → {S[K[x_]][y_][z_] → x[y[z]], K[x_][y_] → x, S[K][x_][y_] → y,
  S[x_][y_][z_] → x[z][y[z]], 0.959}
extra3 → {S[K[x_]][y_][z_] → x[y[z]], K[K][x_][y_][z_] → y, S[K][x_][y_] → y,
  S[x_][y_][z_] → x[z][y[z]], K[x_][y_] → x, 0.79}
extra4 → {S[K][x_][y_] → y, K[K][x_][y_][z_] → y, K[K][x_][y_][z_] → x,
  K[x_][y_] → x, S[K[x_]][y_][z_] → x[y[z]], S[x_][y_][z_] → x[z][y[z]], 0.765625}
```

Interestingly, when looking at all possible rule sets of a given size, the fastest always has the same elements as the rule set one size smaller, but with one new rule added, and in most cases in a different order. Also, the one extra rule that is in all of these fastest rule selections is  $S[K][x\_][y\_]\rightarrow y$ , which again is equivalent to the identity, or “I” combinator  $I[x\_]\rightarrow x$ .

In the graphs in Figure 1, the  $x$  axis represents rules in the order that they show up using the Permutations function. The  $y$  axis shows the average time in seconds for evaluating sets of large combinators.

When one of each extra rule was added to extra4, the rule with a fifth extra rule that was fastest was

```
{S[K][x1_][x2_] → x2, K[K][x1_][x2_][x3_] → x2, S[S][x1_][K][x2_] → x2,
  K[K[x1_]][x2_][x3_] → x1, K[x_][y_] → x, S[K[x1_]][x2_][x3_] → x1[x2[x3]], S[x_][y_][z_] →
  x[z][y[z]]},
```

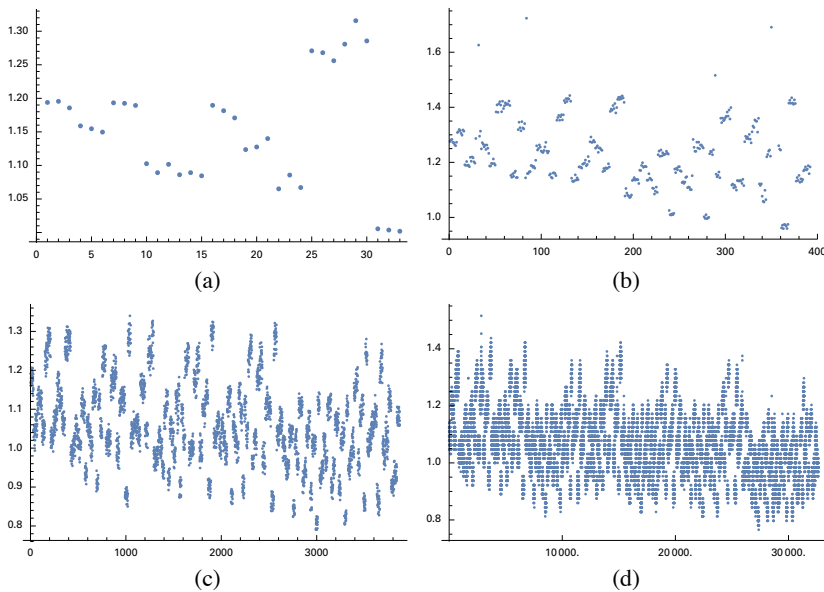
which took 0.915 seconds to go through all the combinators. This rule set was slower than extra4.

To speed up finding rules of size 7, only rules similar to the fastest rule of size 6, but with one new rule part added, were tested.

The fastest of these, however, was slower than the fastest rule set that has one less rule. The rule was

```
{S[K][x1_][x2_] → x2, K[K][x1_][x2_][x3_] → x2, S[S][x1_][K][x2_] → x2,
  K[K[x1_]][x2_][x3_] → x1, K[x_][y_] → x, S[K[x1_]][x2_][x3_] → x1[x2[x3]], S[x_][y_][z_] →
  x[z][y[z]]},
```

which took 0.915 seconds to go through all the combinators. This could be due to the fact that extra rules will generally significantly slow down evolution, if the rule has no effect of skipping steps in evaluation.



**Figure 1.** Timing versus permutation index of implementing combinator rules. This shows how different rules speed up or slow down the evaluation of large combinator rules. The  $x$  axis represents rules in the order that they show up using the Permutations function. The  $y$  axis shows the average time in seconds for evaluating sets of large combinator rules. (a) An extra optimization rule is added in different sets of ways to the standard S and K rules. (b) Two extra rules are added. (c) Three extra rules are added. (d) Four extra rules are added.

Thus, extra rules can only help in the case of large combinator rules, and it appears that only up to a certain point do extra rules continue to speed up evolution. Taking into consideration that extra rules had a modest effect on speeding up evolution of size 100 combinator rules, and that extra rules would slow down the evolution of smaller combinator rules, in the end only the combinator rule that corresponds to the “I” combinator was added to the actual EvaluateCombinator function.

#### 4. Random Long-Running Combinators

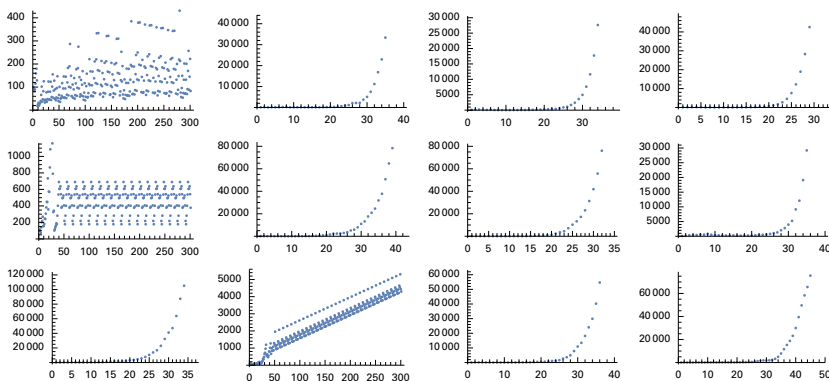
The way in which large combinator rules were evaluated was investigated in order to learn more about why the rules found only had a modest effect on the speeding up of combinator evolution.

First, the evolution of 400 randomly generated size 100 combinator rules was run, and most of these combinator rules quickly did one of two things. They either began growing exponentially or began having repetitive behavior. Only one, in fact, appeared to be doing neither.

First of all, it is known that two of the 16 896 possible size 7 combinators do not reach fixed points, and these combinators grow exponentially [2]. Also, “At size 8, out of all 109 824 combinator expressions it appears that 49 show exponential growth. And many more show roughly linear growth.” [2] It seems that this becomes a trend, with even larger combinators being more and more likely to grow at an exponential rate, at least for a large number of steps. This could be partially due to a larger combinator being more likely to contain a smaller combinator part that grows exponentially.

The behavior of larger combinators that do not halt before reaching certain computational limits was investigated. Four hundred random size 100 SK combinators were generated, using `SeedRandom[1]` as a starting seed for the random combinator generator. These combinators were then evaluated, with evolution pausing when one of three outcomes happened: (1) the combinator evolution reached a fixed point (no longer changed with replacement rules applied); (2) the combinator evolution reached 300 steps; or (3) the size of the combinator went above 2000.

Combinators that led to criterion 1 (reached a fixed point) were filtered out, leaving only combinators that took many steps and/or became very large without halting. Only 66 of the original 400 randomly generated combinators had not halted or had grown too large by this number of steps. These 66 combinators were then evaluated a second time, again for a maximum of 300 steps, but for a maximum combinator size of 200 000, to see if they would halt after reaching this size. Figure 2 shows combinators that did not halt after 300 steps or before reaching size 200 000.



**Figure 2.** (*continues*)

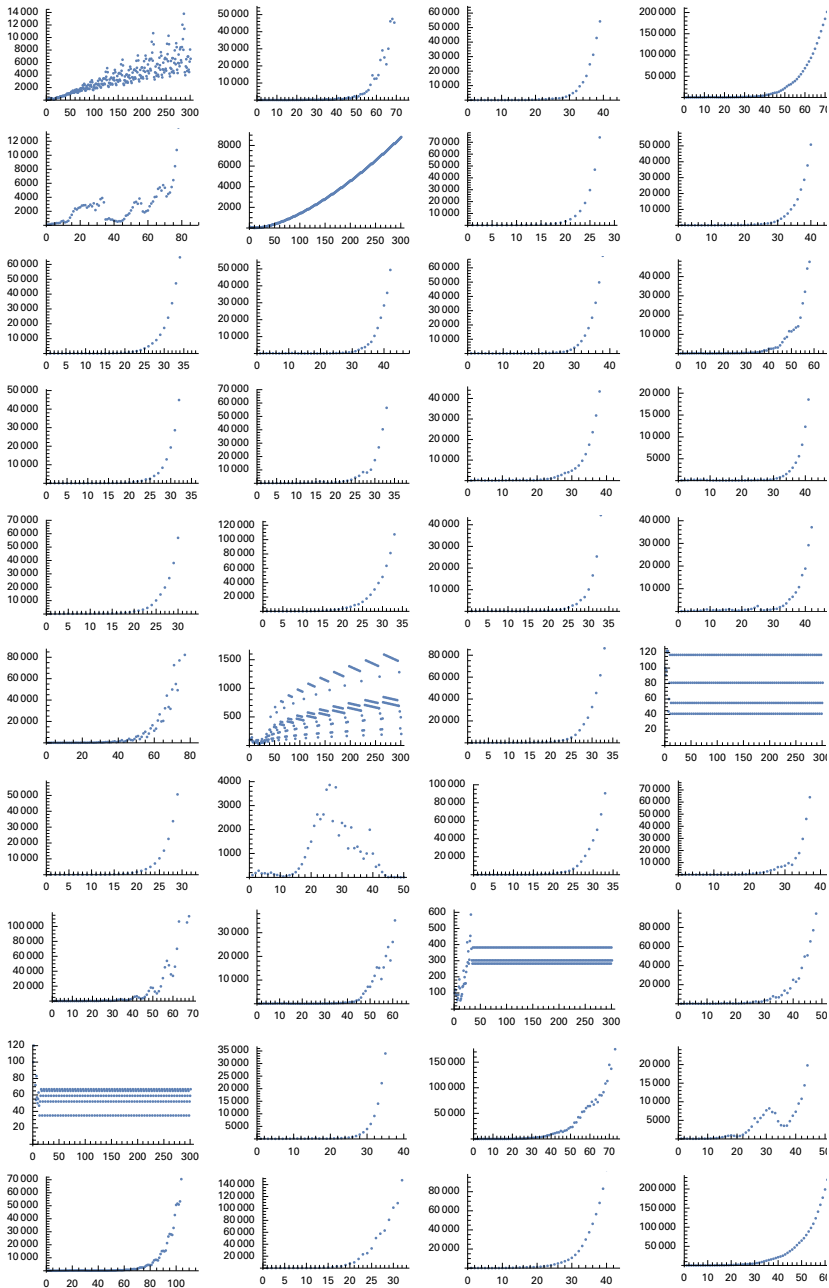
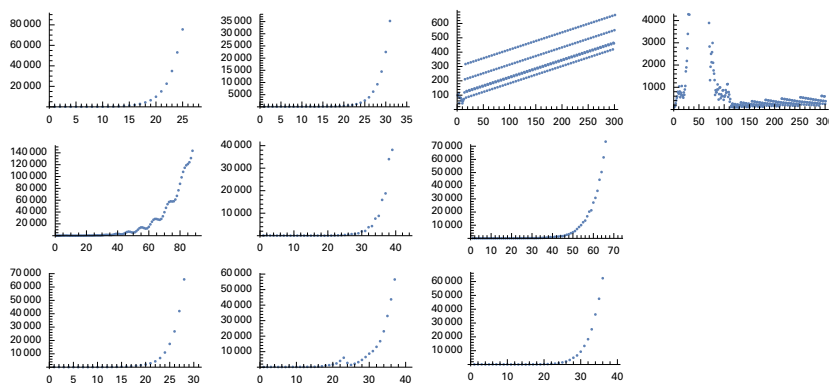


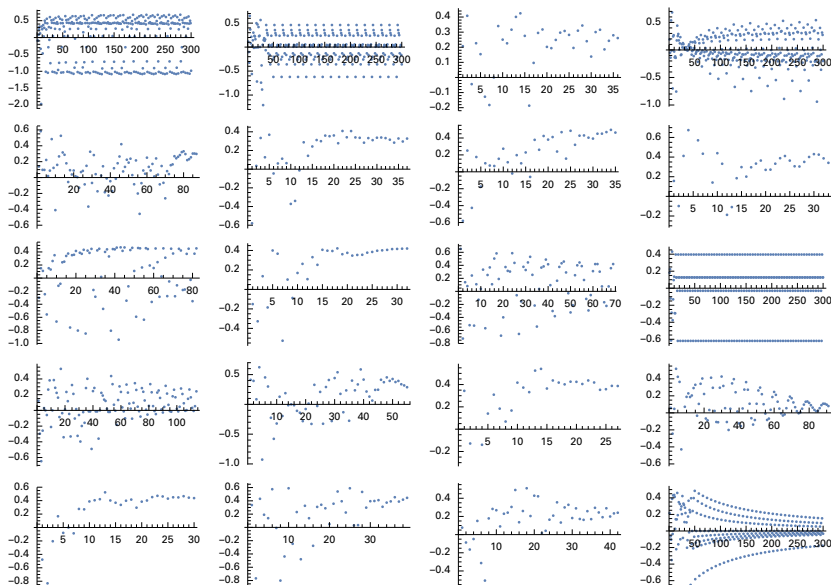
Figure 2. (continues)





**Figure 2.** Randomly selected non-terminating combinator growth. This shows the growth of randomly selected combinators that were also selected under the criterion that they not terminate after many steps.

Figure 3 shows the same plots, but with the natural logarithm taken of each size point, and then with the differences taken between those. It shows that most of the combinators that do grow exponentially still appear to be exhibiting complex behavior.



**Figure 3.** (*continues*)

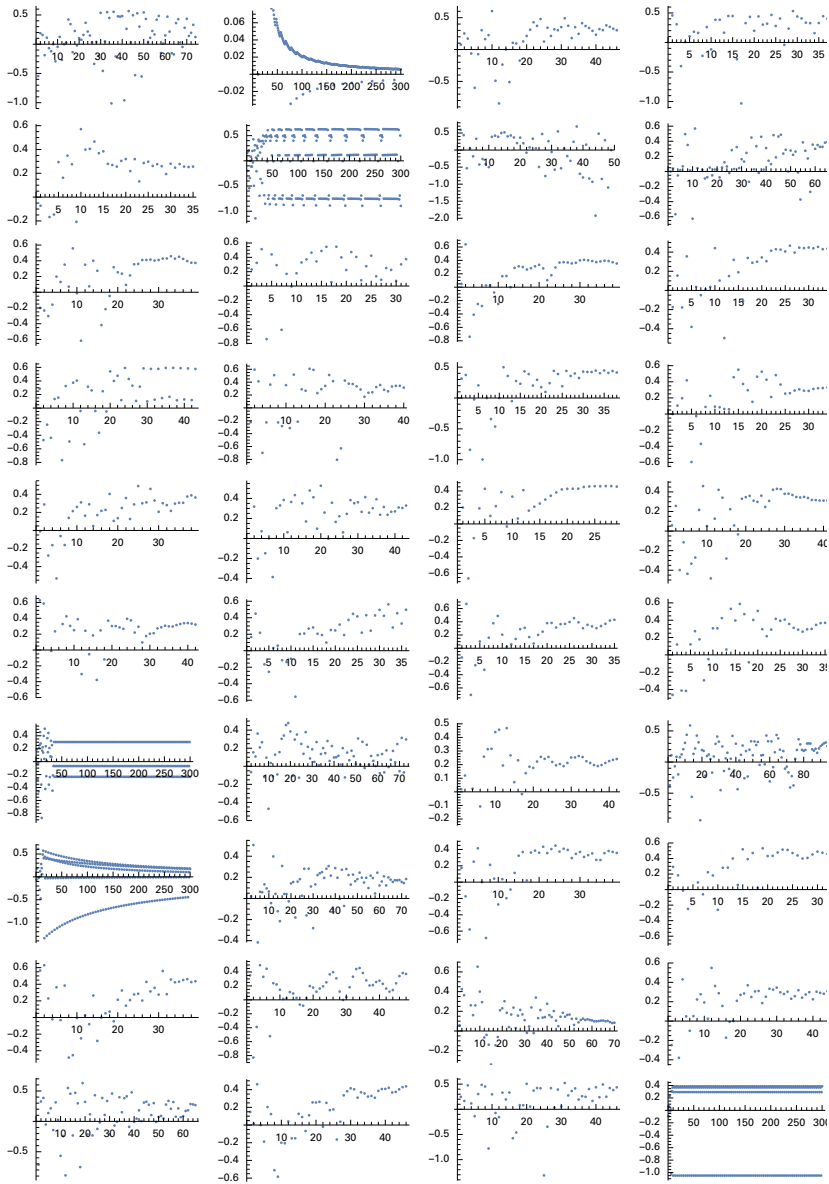
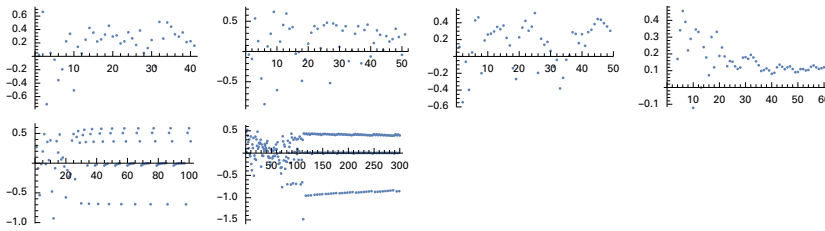


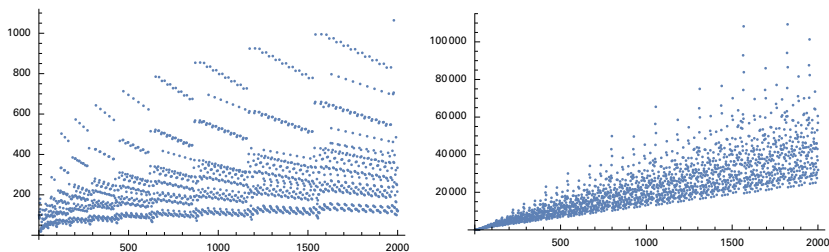
Figure 3. (continues)



**Figure 3.** Logarithmic differences of randomly selected non-terminating combinator growth. This shows the growth of randomly selected combinators that do not terminate after many steps. The y axis shows the logarithms of the differences between combinator lengths from step to step.

Of the 66 combinators, 53 appear to have been selected by going above length 2000 before they reached step 200. One of them got above size 2000 but below 200 000 (it reached a maximum size of 3861), but reached a fixed point anyway before step 300. Four of them got to step 300 by beginning to loop through the same few values repeatedly. Six of them began to grow in what appears to be a repetitive pattern.

Figure 4 shows the two combinators that appeared to be showing complex behavior without growing exponentially, this time with a cut-off of 2000 steps. They both appeared to show complicated behavior without becoming repetitive or growing exponentially, although the behavior may be nested.



**Figure 4.** Two large combinators that appear to show complicated behavior without becoming repetitive or growing exponentially.

It should be noted, however, that the combinators that appeared to be growing at an exponential rate or that seemed to be exhibiting complex behavior even after many steps or reaching a large size could still potentially reach a fixed point if run for more steps or to larger sizes. However, with the limited computational resources given, this was not observed in most of the large combinators.

## 5. Random Terminating Combinators

Figure 5 shows 62 combinators selected from 200 randomly chosen size 1000 combinators that did terminate. All of the ones that reached a fixed point did so within 100 steps (one ran for exactly 100 steps). The largest size any of them reached at any time during its evaluation was 154 820 (it ran for a total of 44 steps).

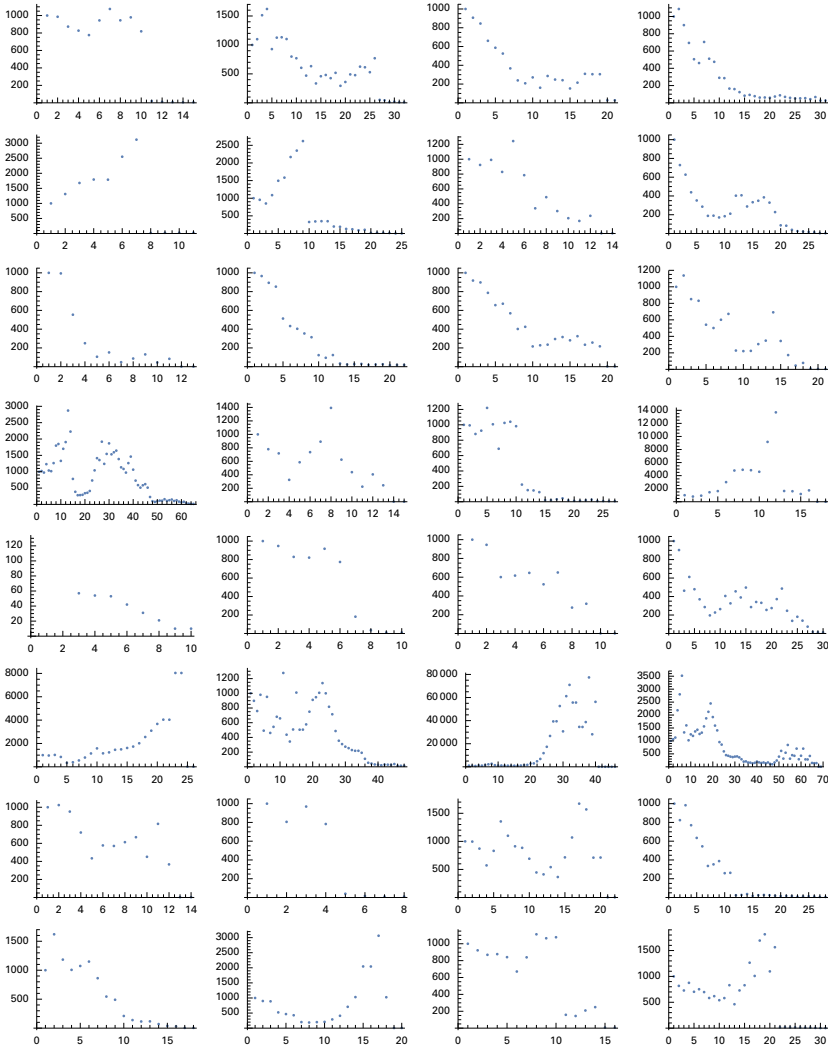
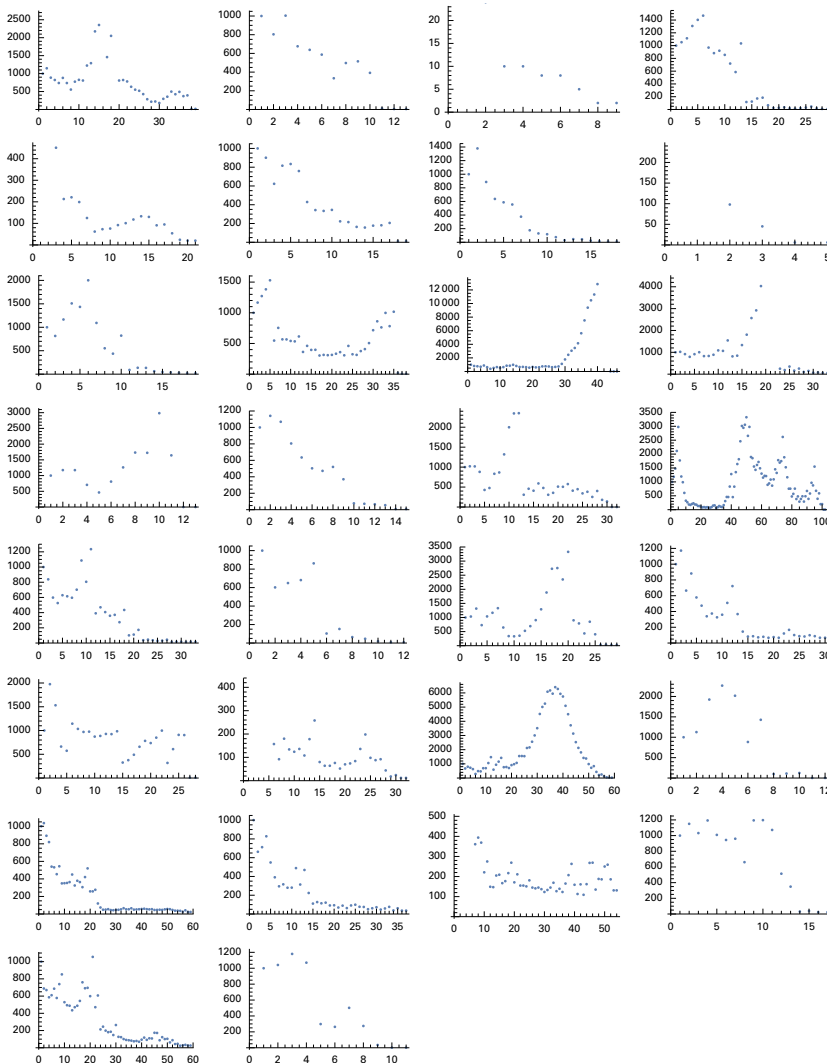


Figure 5. (*continues*)



**Figure 5.** Randomly selected terminating combinator growth. This shows the growth of randomly selected combinators that terminate before a given large number of steps.

## 6. Frequency of Optimization Rule Usages

This shows what fraction of the time an extra optimization rule (besides the usual S and K rules) was used during the attempt at optimized evolution. The combinators are sorted by how often the rules were used. Each letter corresponds to a different size 1000 combinator.

{A → 0.0257056, B → 0.024629, C → 0.0246223, D → 0.024045, E → 0.0235916,  
 F → 0.0225295, G → 0.0222673, H → 0.021963, I → 0.0212766, J → 0.0210058,  
 K → 0.0209823, L → 0.0209536, M → 0.0209166, N → 0.0208107, O → 0.0206622,  
 P → 0.0205735, Q → 0.0205672, R → 0.0204877, S → 0.0204736, T → 0.0204562,  
 U → 0.0202132, V → 0.0200748, W → 0.0200241, X → 0.0199806, Y → 0.0196201,  
 Z → 0.0193757, A1 → 0.0192907, B1 → 0.0192356, C1 → 0.0191463,  
 D1 → 0.0190822, E1 → 0.0189619, F1 → 0.0188708, G1 → 0.0188185,  
 H1 → 0.0185733, I1 → 0.0185015, J1 → 0.0184146, K1 → 0.0183781,  
 L1 → 0.0183239, M1 → 0.0182556, N1 → 0.018183, O1 → 0.0181257,  
 P1 → 0.0180765, Q1 → 0.0179722, R1 → 0.0178833, S1 → 0.0173198,  
 T1 → 0.0171956, U1 → 0.0161381, V1 → 0.0160847, W1 → 0.0159823,  
 X1 → 0.015799, Y1 → 0.0156495, Z1 → 0.0155909, A2 → 0.0155469, B2 → 0.015438,  
 C2 → 0.0148621, D2 → 0.0146095, E2 → 0.0141228, F2 → 0.0138126,  
 G2 → 0.0107118, H2 → 0.00715231, I2 → 0.00564804, J2 → 0.0050268}

Figure 6 labels the combinator optimization rules. Figure 7 breaks down how often each of these labeled rules is used individually in evaluation out of 100 000 and also shows the number of times out of 100 000 when a part of the combinator evolution does not match any rule.

```

a → S[S[x1_][K][x2_] → x2
b → S[S[K][x1_][x2_] → x1[x2]
c → S[S][K][x1_][x2_] → x1[x2][x1]
d → K[K][x1_][x2_][x3_] → x2
e → K[K[x1_][x2_][x3_] → x1
f → K[K[x1_][x2_][x3_] → x1
g → S[x1_][K[x2_][x3_] → x1[x3][x2]
h → S[K[x1_][x2_][x3_] → x1[x2][x3]
i → S[S][x1_][x2_][x3_] → x2[x3][x1[x2][x3]]
j → S[S[x1_][x2_][x3_] → x1[x2[x3]][x3[x2[x3]]]
k → S[K][x1_][x2_] → x2
l → K[x_][y_] → x
m → S[x_][y_][z_] → x[z][y[z]]
n → x_ → x

```

**Figure 6.** Labeling optimization rules for the table in Figure 7.

	a	b	c	d	e	f	g	h	i	j	k	l	m	n
A	45	93	405	546	45	11	270	138	546	186	286	1823	808	94798
B	213	225	6	599	33	18	198	189	362	295	325	2484	888	94165
C	51	87	29	298	414	15	327	247	283	189	523	1816	1148	94574
D	12	257	17	82	29	48	149	596	125	580	510	1485	1489	94621
E	76	76	84	353	260	126	294	134	269	118	571	2359	1192	94089
F	24	48	83	238	191	83	274	215	215	453	429	2134	1359	94254
G	78	78	90	90	66	84	394	149	310	173	716	2298	1128	94347
H	30	76	45	197	121	76	288	257	151	273	682	2196	1621	93987
I	89	113	122	215	85	122	300	178	349	150	405	2610	1325	93937
J	21	106	38	458	168	14	322	147	192	144	493	1796	1020	95084
K	5	21	72	257	370	26	257	77	216	108	689	2294	1353	94256
L	64	64	86	278	107	21	257	150	214	214	641	2651	1219	94035
M	7	145	51	232	174	138	254	240	203	196	450	2142	1118	94647
N	22	185	115	393	153	70	118	80	221	211	511	2986	1132	93802
O	85	109	193	169	97	133	169	230	290	205	387	2127	1112	94696

**Figure 7.** (continues)

	a	b	c	d	e	f	g	h	i	j	k	l	m	n
P	16	43	155	187	224	32	197	155	352	165	533	2473	911	94558
Q	27	127	63	136	399	136	263	109	181	190	426	2600	870	94473
R	31	13	82	308	195	258	145	176	201	44	597	2709	1276	93967
S	25	99	49	345	173	74	247	222	123	271	419	2935	1184	93833
T	26	22	26	232	318	45	310	90	303	131	542	1421	1144	95389
U	61	104	141	202	37	37	147	202	392	233	466	3001	1133	93844
V	25	157	96	298	253	40	243	162	233	172	329	2559	895	94539
W	28	51	79	222	88	97	357	121	334	167	459	2647	973	94377
X	6	11	91	194	194	11	360	337	280	6	508	2546	1222	94234
Y	37	31	112	187	156	93	187	212	336	137	473	1682	1046	95310
Z	20	20	173	81	243	66	173	234	228	141	557	1071	1135	95856
A1	12	47	39	304	126	20	201	158	205	233	584	1333	1325	95412
B1	25	25	50	225	250	75	200	200	250	125	500	2698	999	94379
C1	10	105	33	119	95	43	282	234	115	363	516	1289	1289	95507
D1	47	102	37	168	140	84	279	112	205	307	428	2523	1378	94192
E1	28	42	83	249	194	69	180	125	277	138	512	2920	955	94228
F1	5	118	15	159	67	62	257	293	129	221	560	1568	1203	95341
G1	39	165	24	194	204	53	179	58	189	247	529	2454	1174	94490
H1	0	6	0	296	661	3	340	29	7	498	17	606	890	96647
I1	85	23	85	128	124	43	113	346	283	274	345	2323	1444	94382
J1	46	63	11	193	151	7	319	260	200	119	474	2483	1105	94570
K1	5	49	16	65	139	84	272	383	92	117	614	1283	1188	95691
L1	43	36	25	140	36	36	428	122	263	324	378	1120	1688	95360
M1	16	54	13	432	115	51	192	67	192	144	550	2513	1148	94514
N1	86	33	7	178	119	26	198	224	250	79	619	1779	1252	95151
O1	45	58	58	135	103	122	231	180	212	96	572	2410	1202	94575
P1	7	31	50	474	67	22	87	50	87	132	800	1250	1382	95560
Q1	54	34	41	183	332	27	176	122	264	109	454	2197	943	95063
R1	9	48	18	254	169	16	229	338	64	105	537	1301	1249	95662
S1	127	80	27	107	67	160	74	134	201	154	602	2213	1237	94817
T1	6	18	54	223	235	127	97	157	133	121	549	2552	887	94841
U1	7	22	31	109	64	22	107	97	191	164	800	814	1113	96459
V1	75	0	19	115	71	14	105	104	441	113	551	1398	909	96084
W1	15	29	18	147	50	24	144	162	197	303	509	2028	1266	95108
X1	5	25	5	231	85	125	231	181	201	85	406	2112	828	95481
Y1	41	36	51	87	31	5	273	264	351	62	364	1716	869	95851
Z1	4	31	13	297	66	17	301	201	231	109	288	2328	926	95187
A2	20	55	25	325	115	90	170	195	115	200	245	2050	915	95481
B2	6	9	209	91	49	4	82	147	220	483	242	743	911	96802
C2	63	33	37	187	53	57	240	224	180	83	327	1383	698	96433
D2	24	77	19	89	228	56	100	81	212	121	453	1664	873	96002
E2	5	103	18	166	27	10	161	203	206	127	388	910	1153	96524
F2	15	15	10	175	113	87	277	252	67	62	308	1951	703	95964
G2	1	56	33	151	82	68	173	79	70	91	268	935	653	97341
H2	11	21	35	80	17	2	106	125	95	37	186	544	403	98338
I2	1	23	4	14	9	4	32	38	15	38	387	291	674	98470
J2	20	18	14	28	18	1	25	72	59	57	191	374	606	98518

**Figure 7.** Shows the number of times a rule was used during the evolution of different large combinator. Rules are column labels and combinator are row labels. They are ordered by how many times total a rule (besides the S and K rules) matched, from most to least. S and K rules are third and fourth from the right, and the rightmost column shows the number of subpatterns of combinator that did not match any combinator rule.

These are the rounded number of times a rule was used on average out of 100 000.

{a  $\rightarrow$  34, b  $\rightarrow$  65, c  $\rightarrow$  60, d  $\rightarrow$  215, e  $\rightarrow$  146, f  $\rightarrow$  58, g  $\rightarrow$  215,  
h  $\rightarrow$  178, i  $\rightarrow$  216, j  $\rightarrow$  181, k  $\rightarrow$  467, l  $\rightarrow$  1918, m  $\rightarrow$  1083, n  $\rightarrow$  95 163}

## 7. Conclusion

Out of the numerous possible rules up to size five that could be made to skip combinator evolution, only 11 were not redundant in some way. This made it easier to look through different possible ways to use them for speeding up combinator evolution. In the end, adding too many extra rules also generally slows down evolution, as each rule needs to be compared with all parts and subparts of a combinator expression.

Optimization rules are generally most helpful for combinators that terminate. As for non-terminating combinators, we would most likely want to look at the evolution step by step, and optimization rules would skip steps in a generally hard-to-predict fashion. The majority of large combinators, say of size 100, appear to not terminate. Most begin to grow exponentially rather quickly, and some start exhibiting repetitive behavior. A couple of interesting cases were found where there was no exponential growth, but there also seemed to be less repetitive behavior, and it is unclear whether these will in fact terminate at some point.

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