



































2. Compare the  $n^2$  entries  $A_{ij}$  and  $B_{ij}$ . This step requires at most  $n^2$  comparisons. (This maximal value is needed in the case that  $f$  is a fixed point system.)
3.  $f$  is a fixed point system if and only if the matrices  $A$  and  $B$  are equal.

It is well known that matrix multiplication requires  $2n^3 - n^2$  addition or multiplication operations. Since  $t + 1 < n \log_2(q) + 2$ , the number of operations required in step 1 is bounded above by  $(2n^3 - n^2)(n \log_2(q) + 2)$ . Summarizing, we have the following upper bound  $N(n, q)$  for the number of operations in steps 1 and 2:

$$N(n, q) := (2n^3 - n^2)(n \log_2(q) + 2) + n^2.$$

For a fixed size  $q$  of the finite field  $F_q$  used it holds that

$$\lim_{n \rightarrow \infty} \frac{N(n, q)}{n^4} = 2 \log_2(q)$$

and we can conclude  $N(n, q) \in O(n^4)$  for a fixed  $q$ . The asymptotic behavior for a growing number of variables and growing number of field elements is described by

$$\lim_{\substack{n \rightarrow \infty \\ q \rightarrow \infty}} \frac{N(n, q)}{n^4 \log_2(q)} = 2.$$

Thus,  $N(n, q) \in O(n^4 \log_2(q))$  for  $n, q \rightarrow \infty$ .

It is necessary to comment on the arithmetic operations performed during the matrix multiplications. Since the matrices are elements of the matrix monoid  $M(n \times n; E_q)$ , the arithmetic operations are operations in the monoid  $E_q$ . The addition (resp., multiplication) operation on  $E_q$  requires an integer number addition (resp., multiplication) and a reduction using  $\text{red}_q$  (see Section 2). For a detailed description of integer number representation and arithmetic in typical computer algebra systems see [28, Chapter 4].

The reduction  $\text{red}_q(a)$  of an integer number  $a \in \mathbb{N}_0$ ,  $a \geq q$  is obtained as the degree of the remainder of the polynomial division  $\tau^a \div (\tau^q - \tau)$ . According to [28, Section 4.6.5] this division requires  $O(2(\deg(\tau^a) - \deg(\tau^q - \tau))) = O(2(a - q))$  integer number operations. However, we know that the reductions  $\text{red}_q(\cdot)$  are applied to the result of (regular integer) addition or multiplication of elements of  $E_q$  and therefore

$$a - q \leq \begin{cases} 2(q - 1) - q = q - 2 \\ (q - 1)^2 - q = q^2 - q + 1. \end{cases}$$

As a consequence, in the worst case scenario, one addition (resp., multiplication) in the monoid  $E_q$  requires  $O(q)$  (resp.,  $O(q^2)$ ) regular integer number operations.

Since  $E_q$  is a finite set and only the results of  $n^2$  pairwise additions and  $n^2$  pairwise multiplications are needed, the numbers are stored in a table after the first time they are calculated while the algorithm is running. Since most of the commercial and freely available computer algebra systems provide implementations of finite field arithmetic and arithmetic of polynomials over a finite ring, the implementation of our algorithm is a very simple task. We have successfully implemented the algorithm in Maple<sup>TM</sup>. The code can be made available from the author upon request.

Since the matrix multiplications dominate the complexity of the algorithm, for very large systems more efficient matrix multiplication algorithms could be used. Indeed, using Strassen's algorithm [29, 30], a complexity of  $O(n^{\omega+1} \log_2(q))$ , where  $\omega \leq \log_2(7) \approx 2.807$  could be reached. The performance could be substantially improved through the use of parallelization techniques [31, 32].

It is pertinent to mention that while this article was being peer-reviewed, the article [14] was published, in which an algorithm of complexity  $O(n^3 \log_2(n \log_2(q)))$  is presented, that is used to determine whether an  $n$ -dimensional linear dynamical system  $L : R^n \rightarrow R^n$  over a finite ring  $R$  of cardinality  $q = |R|$  is a fixed point system. While the authors of [14] did perform a complexity analysis of their algorithm, they did not elaborate on the details of a concrete computer implementation of the arithmetic operations on the finite ring  $R$  and thus, did not provide information on the computational cost of such operations.

In [13] it is stated that in order to exploit their results algorithmically, an algorithm would be required to determine whether or not a linear dynamical system  $L : R^n \rightarrow R^n$  over a finite ring  $R$  is a fixed point system. The algorithm presented in [14] provides this missing ingredient. Nevertheless, for a given  $f \in M F_n^n(\mathbb{F}_q)$ , prior to being able to apply the algorithm from [14], a linear dynamical system  $L(f) : (\mathbb{Z} / (q - 1))^n \rightarrow (\mathbb{Z} / (q - 1))^n$  and a Boolean monomial dynamical system  $T(f) \in M F_n^n(\mathbb{F}_2)$  need to be constructed (see [13] for the details). The overall computational complexity of an algorithm exploiting the results of [13] and [14] still remains to be investigated.

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